

Konzepte objektorientierter Programmierung

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Exercises 8: Aliasing

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Java Modeling Language JML

- Available at:
<http://jmlspecs.org/>
- Gary T. Leavens, Yoonsik Cheon.
Design by Contract with JML.
<ftp://ftp.cs.iastate.edu/pub/leavens/JML/jmldbc.pdf>
- “The Java Modeling Language (JML) is a behavioral interface specification language that can be used to specify the behavior of Java modules.”



JML Notation

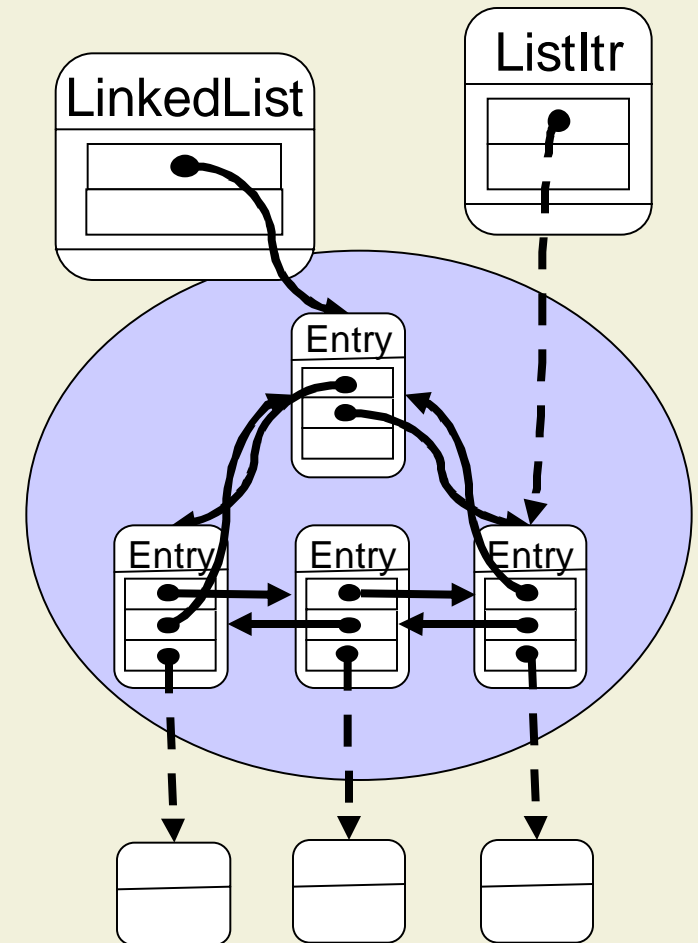
- Additional stuff in special comments:
`//@ till end of line`

```
/*@ requires x > 0;  
   @ ensures x == 0;  
   @* /
```

- In the lecture examples, we sometimes leave out the comments and sometimes the @
- More later

Alias Control by Extended Typing

- We introduce different types for the different roles of objects
 - **peer** types for objects in the **same context as this** (interface objects)
 - **rep** types for representation objects in the **context owned by this**
 - **readonly** types for argument objects **in any context**
- Type rules replace the programming discipline



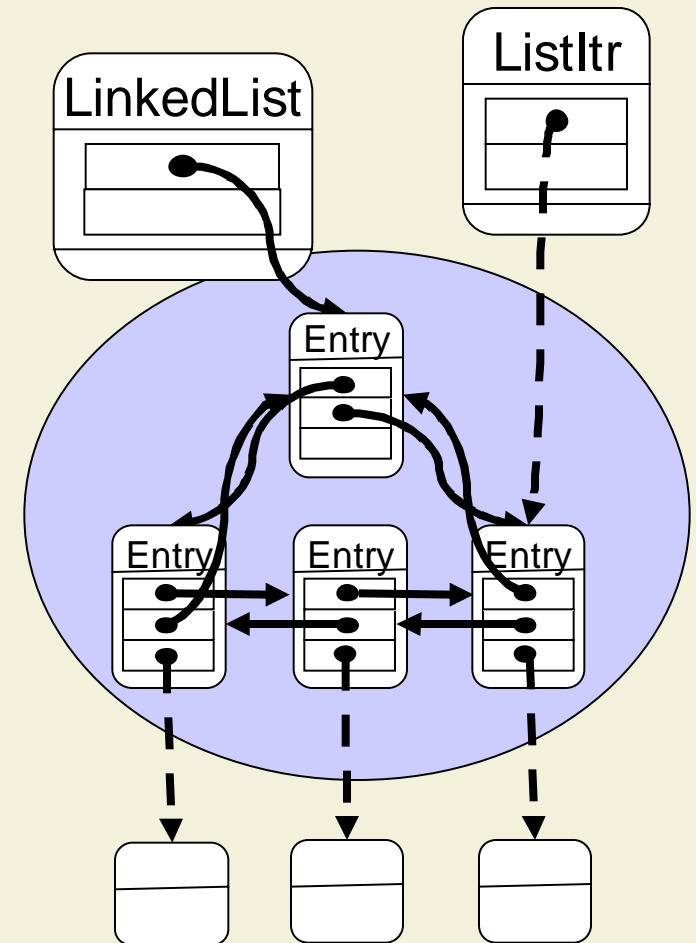
Linked List Example

```
public class LinkedList {  
    public void add( readonly Object o ) {...}  
  
    public readonly Object getFirst() {...}  
    ...  
}
```

```
LinkedList l = new peer LinkedList();  
l.add( new peer Integer(3) );  
l.add( new peer Object() );  
  
peer Object o = (peer Object) l.getFirst();
```

Alias Control by Extended Typing

- We introduce different types for the different roles of objects
 - peer types for objects in the **same context as this** (interface objects)
 - Rep types for representation objects in the **context owned by this**
 - Readonly types for argument objects **in any context**
- Type rules replace the programming discipline



Homework – Exercise 1

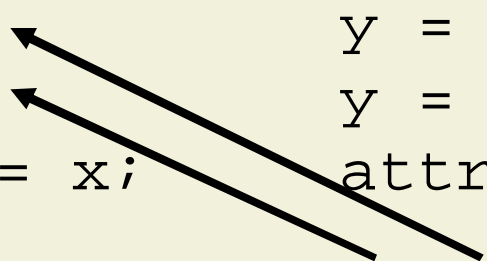
```
class Room {  
    private Human m;  
    void enter( Human m ) {}  
    Human leave() {}  
}  
public class Building {  
    private Room[] rooms;  
    void enter( Human m ) {  
        rooms[0].enter( m ); }  
    Human leave() {  
        rooms[0].leave( m ); }  
}
```

Homework – Exercise 1

```
class Room {  
    private /*arg*/ Human m;  
    void enter(/*arg*/ Human m ) {}  
    /*arg*/ Human leave() {}  
}  
  
public class Building {  
    private /*rep*/ Room[] rooms;  
    void enter( /*arg*/ Human m ) {  
        rooms[0].enter( m ); }  
    /*arg*/ Human leave() {  
        rooms[0].leave( m ); }  
}
```


Homework – Exercise 2

```
public class Assignments {  
    private /* rep */ Integer attr_x;  
    private /* arg */ Integer attr_y;  
    private Integer attr_z;  
    public void m( /* arg */ Integer a ) {  
        /* rep */ Integer x = new /* rep */  
            Integer(4);  
        /* arg */ Integer y = new Integer(3);  
        Integer z = new Integer(2);  
  
        x = y;          y = x;  
        x = z;          y = z;  
        attr_x = x;     attr_y = a;  
    }  
}
```



Error!

(Simplified) Programming Discipline

■ Rule 1: No Role Confusion

- Expression with one alias mode must not be assigned to variables with another mode

■ Rule 2: No Representation Exposure

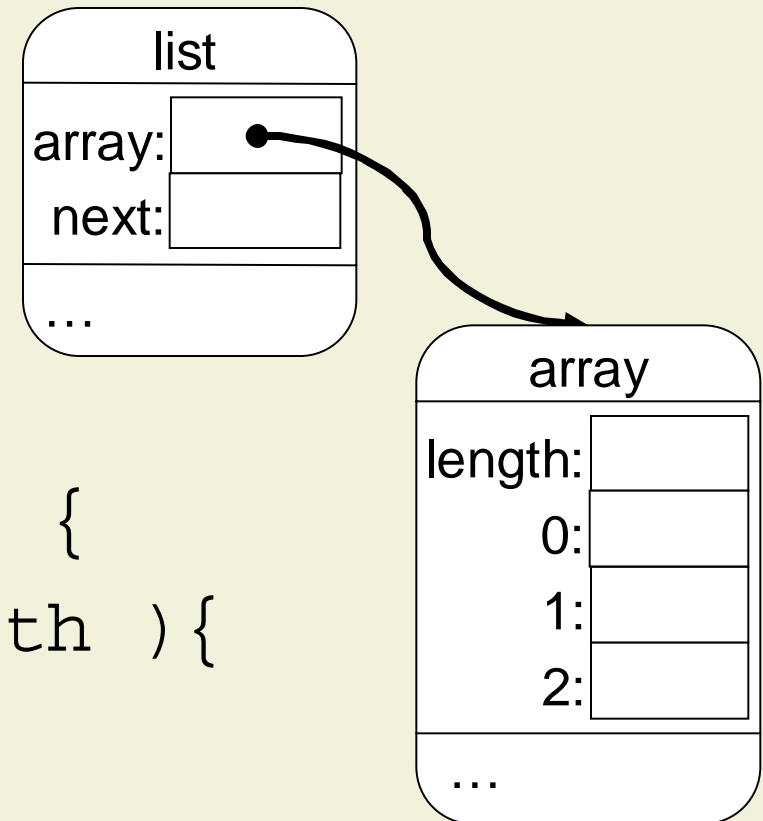
- rep-mode must not occur in an object's interface
- Methods must not take or return rep-objects
- Fields with rep-mode may only be accessed on **this**

■ Rule 3: No Argument Dependence

- Implementations must not depend on the state of argument objects

ArrayList

```
class ArrayList {  
    protected int[] array;  
    protected int next;  
  
    public void add(int i) {  
        if( next==array.length ) {  
            resize( );  
        }  
        array[ next ] = i;  
        next++;  
    }  
}
```



ArrayList

```
public void setElems( int[] ia ) {  
    array = ia;  
}
```

```
public int[] getElems() {  
    return array;  
}
```

- setElems is *Capturing* an external array
- getElems is *Leaking* the internal representation

ArrayList – Main Program

```
public static void main( String[] args ) {  
    int[] myarr = new int[10];  
  
    for(int i=0; i < myarr.length; ++i)  
        myarr[i] = i;  
  
    ArrayList al = new ArrayList();  
  
    al.setElems( myarr );  
  
    myarr[0] = 42;  
}
```

Annotated ArrayList

```
class ArrayList {  
    protected /* rep */ int[] array;  
    protected int next;
```

- Array is part of the internal representation
- Marking it as `/* rep */` makes this explicit for the programmer

ArrayList – setElems

```
public void setElems( int[] ia ) {  
    array = new /* rep */ int[ ia.length ];  
    System.arraycopy  
        (ia, 0, array, 0, ia.length );  
    next = ia.length;  
}
```

- The input array is of default type and can not directly be assigned to the `/* rep */ array`
- We have to create our own copy that can be used as internal representation → no *Capturing*

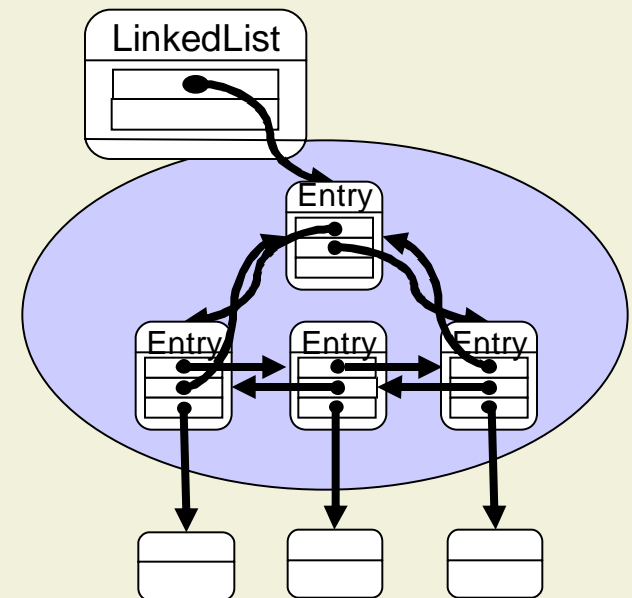
ArrayList – getElems

```
public int[] getElems() {  
    return (int[]) array.clone();  
}
```

- The `/* rep */` array can not be passed out directly as default array
- We create a clone of the internal array and return it as result
- The internal representation is still encapsulated → no *Leaking*

Annotated LinkedList – Setup

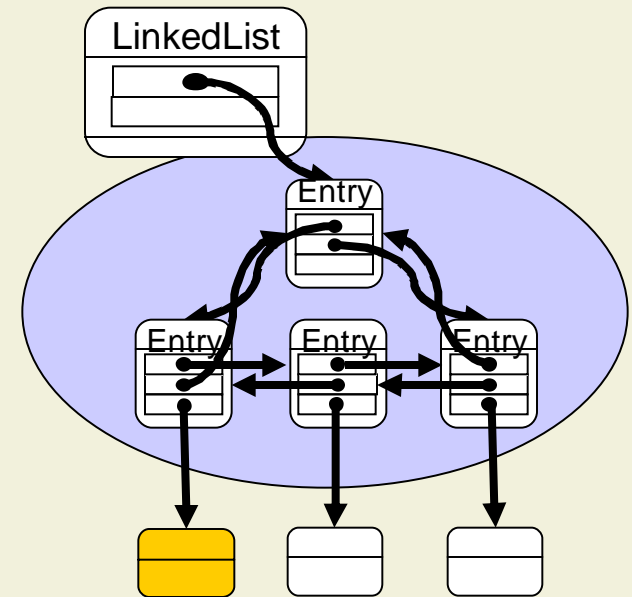
```
class LinkedList {  
    private /* rep */ Entry header;  
    private int size;  
  
    public LinkedList() {  
        header = new /* rep */  
            Entry(null, null, null);  
        header.next = header;  
        header.previous = header;  
  
        size = 0;  
    }  
}
```



LinkedList – add

```
public void add( /* arg */ Object o ) {
    /* rep */ Entry newE =
        new /* rep */ Entry( o, header,
            header.next );
```

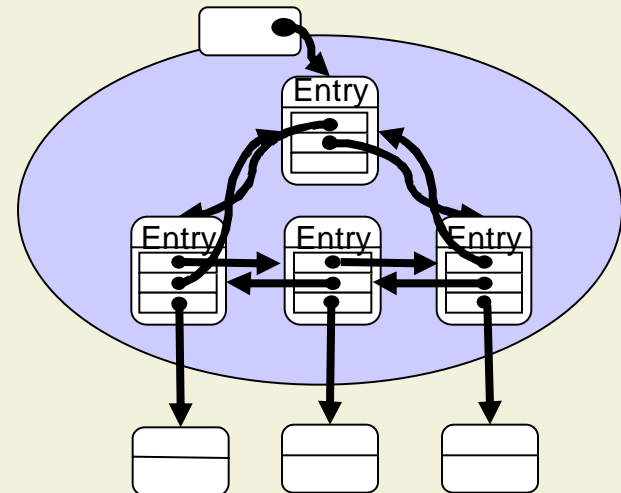
Linked list



```
public /* arg */ Object get( int idx ) { ... }
```

Helper Class Entry

```
class Entry {  
    private /* arg */ Object element;  
    private Entry previous, next;  
  
    public Entry( /* arg */ Object o, Entry p,  
                 Entry n ) {  
        element = o;  
        previous = p;  
        next = n;  
    }  
}
```

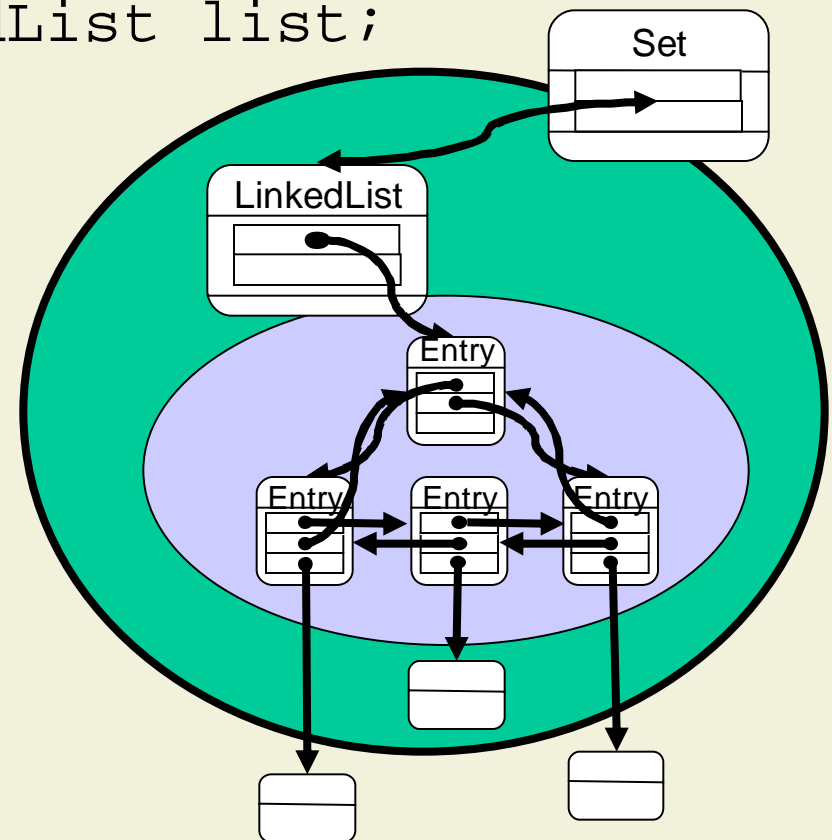


Some Main Program

```
public static void main( String[] args ) {  
  
    LinkedList ll = new LinkedList();  
  
    ll.add( new Integer(20) );  
    ll.add( new Float(2.2f) );  
  
    System.out.println( "Element[0]: " +  
        ll.get(0) );  
    System.out.println( "Element[1]: " +  
        ll.get(1) );  
}
```

Class Set as user of LinkedList

```
public class Set {  
  
    private /* rep */ LinkedList list;  
  
    public Set() {  
        list = new /* rep */  
            LinkedList();  
    }  
  
    ...  
}
```



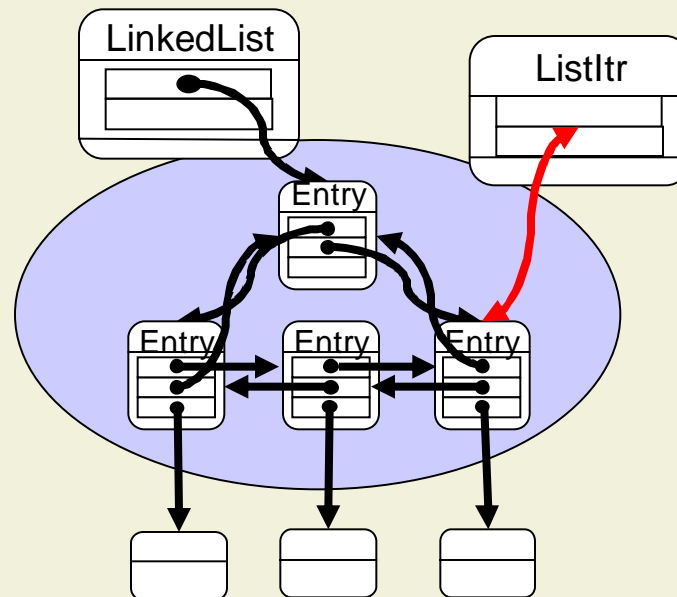
Problems with Annotations

- No help from the compiler, just comments
- Other developers do not have to follow the annotation
- Subclasses do not have to follow the annotation

```
public class BadArrayList  
    extends ArrayList {  
        public int[] leakArray() {  
            return array;  
        }  
    }  
}
```

Just one Object can be the Owner

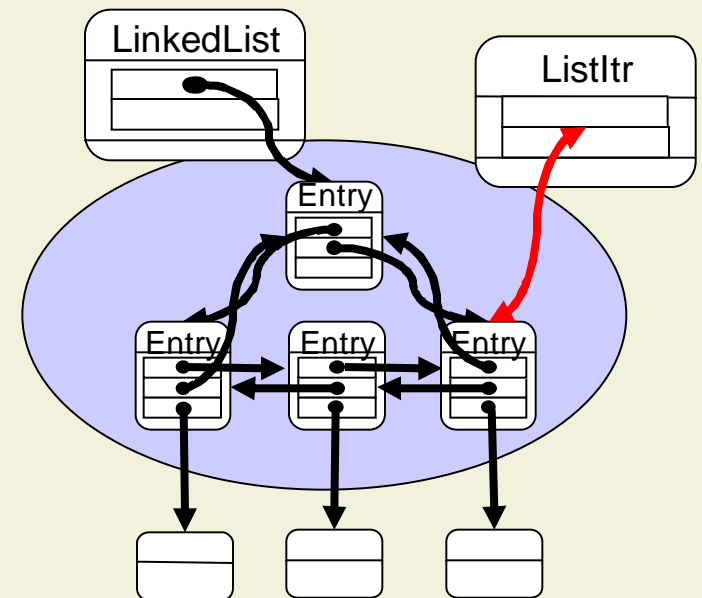
- Sometimes we would like to allow multiple objects access to the internal representation
- E.g. an iterator access to a list



Problems with two Owners

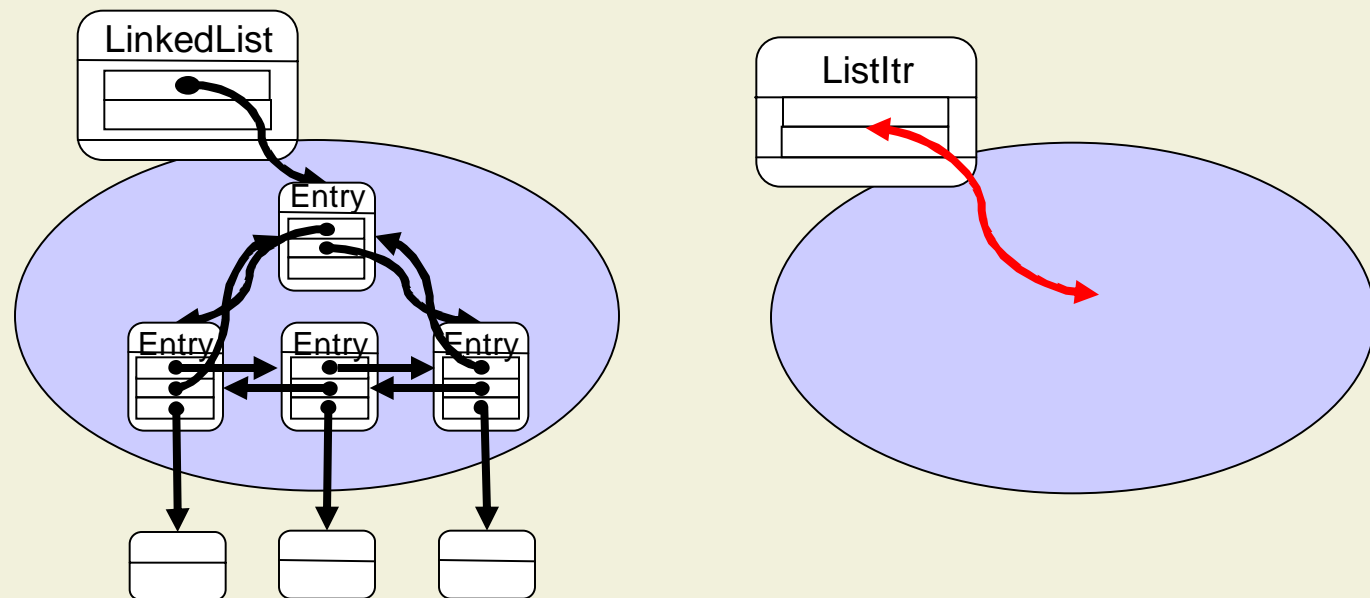
```
ListItr getItr() {  
    ListItr res = new ListItr();  
    res.current = header.next;    }
```

```
class ListItr {  
    /*rep*/ Entry current;  
    /*arg*/ Entry current;  
    Entry current;  
}
```



Problems with two Owners

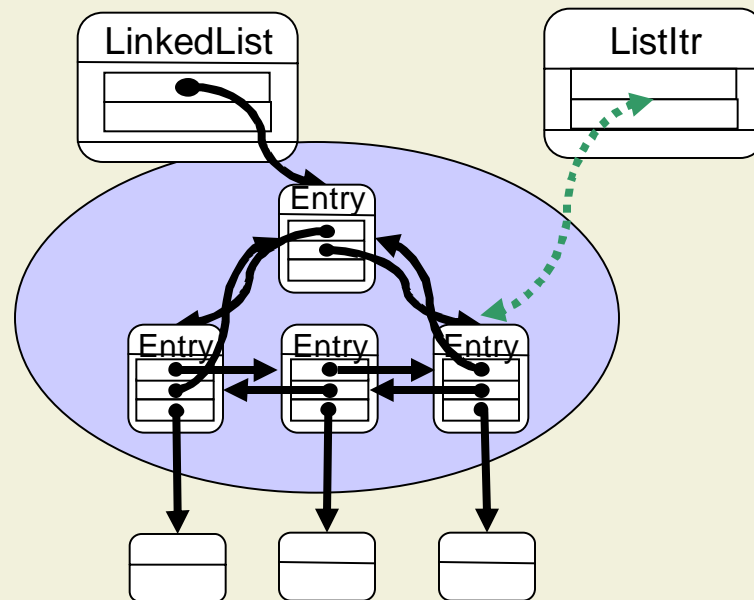
`/* rep */ Entry current;`
would denote an Entry object in Listltr's own context!



Problems with two Owners

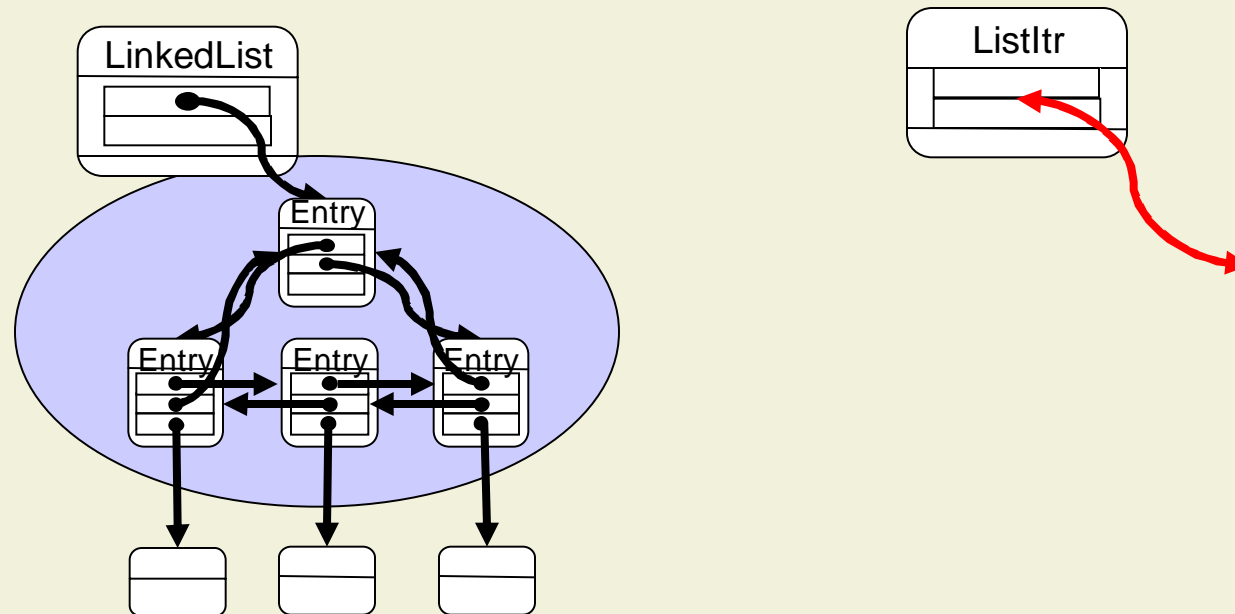
```
/* arg */ Entry current;
```

would denote a readonly reference to an Entry object in the LinkedList context, which might be OK!



Problems with two Owners

Entry current;
would denote an Entry object on Listltr's level!



C++ const

- “When using a pointer, two objects are involved: the pointer itself and the object pointed to. ‘Prefixing’ a declaration of a pointer with `const` makes the object, but not the pointer, a constant. To declare a pointer itself, rather than the object pointed to, to be a constant, we use the declarator operator `*const` instead of plain `*`.”
- Bjarne Stroustrup: “*The C++ Programming Language, Third Edition*”, page 94

C++ const is not transitive

```
class Attr {  
    Attr(int i) {x = i;}  
    int x;  
};
```

```
class A {  
    A() { attr =  
        new Attr(5); }  
    Attr *attr;  
};
```

```
class Main {  
    Main() { a =  
        new A(); }  
};
```

```
    const A *a;  
};
```

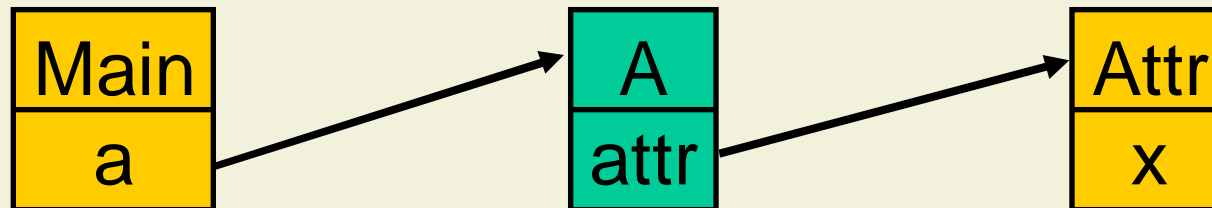
C++ const is not transitive

```
Main *var = new Main();
```

```
var->a->attr->x = 9;
```

```
var->a->attr = new Attr(9); Compilation Error!
```

```
var->a = new A();
```



Readonly Access in Java

- Can be explicitly modeled by using read-only super-interfaces for all classes
- Only return read-only interface to clients
- Problems:
 - Not transitive, we need to change all involved interfaces
 - Hard to reuse existing libraries
 - Not safe, the representation can still leak and it is easy to cast the reference to the read-write type

Programming Discipline

- The programming discipline was just comments for the programmer
- No enforcement by the compiler
- No support for multiple developers
- Easy to make errors even for one developer

Extended Types as a Solution

- We extend the type system of our language to support additional modifiers on types
- Not just annotations in comments, but real changes to the language, **new keywords** and a **different type system**
- This provides us with **direct support from the compiler** in checking whether programs are correct
- We have an **experimental Java compiler** with support for extended types → **not standard Java**

Types

- Each class or interface T **introduces two types**
- **Peer type** $peer(T)$
 - Denoted by T in programs
- **Readonly type** $ro(T)$
 - Denoted by **readonly** T in programs
- For each variable, attribute and parameter we can decide whether it should be read-write or read-only

Subtype Relation

- **Subtyping** among peer and readonly types is **defined as in Java**
 - S extends or implements T
 $\Rightarrow \text{peer}(S) < \text{peer}(T)$
 - S extends or implements T
 $\Rightarrow \text{ro}(S) < \text{ro}(T)$
- **Peer types** are **subtypes of** corresponding **readonly types**
 - $\text{peer}(T) < \text{ro}(T)$

Type Scheme Combinator

- Usually, an access of the form $a.b$ has the type of b
- Now we need to combine the extended types of a and b to get the type of the result

$a * b$	$peer(T)$	$ro(T)$
$peer(S)$	$peer(T)$	$ro(T)$
$ro(S)$	$ro(T)$	$ro(T)$

- By using this combinator readonly becomes transitive

When do we need the combinator?

- When accessing attributes of an object:

```
class A: B b;  
a.b      ➔    result has type a * b
```

- When invoking methods:

```
class A: B m( );  
a.m( )   ➔    result has type a * B
```

Questions?